

Problem Set #9

Description

These problems are related to the material covered in Lectures 16-18.

Instructions: Solve any combination of problems that sum to 100 points. Your solutions are to be written up in latex and submitted as a pdf-file with a filename of the form SurnamePset9.pdf.

Collaboration is permitted/encouraged, but you must identify your collaborators, and any references not listed in the course syllabus. The first to spot each typo/error in the problem sets or lecture notes will receive 1-5 points of extra credit.

Problem 1. Complex multiplication (49 points)

Let $\tau = (1 + \sqrt{-7})/2$. In problem 1 of Problem Set 8 you computed $j(\tau) = -3375$. In problem 2 of Problem Set 7 you proved that the endomorphism ring of the elliptic curve $y^2 = x^3 - 35x - 98$ with j -invariant -3375 is equal to $[1, \tau]$, the maximal order (ring of integers) of $\mathbb{Q}(\sqrt{-7})$. Let us now set $g_2 := -4(-35) = 140$ and $g_3 := -4(-98) = 392$ and work with the isomorphic elliptic curve E/\mathbb{C} defined by

$$y^2 = 4x^3 - g_2x - g_3,$$

which is isomorphic to $y^2 = x^3 - 35x - 98$.

We should note that $g_2([1, \tau])$ and $g_3([1, \tau])$ are not equal to 140 and 392, but there is a lattice L homothetic to $[1, \tau]$ for which $g_2(L) = 140$ and $g_3(L) = 392$ (you computed this lattice L in problem 2 of Problem Set 8). In particular, $\tau L \subseteq L$, thus τ satisfies condition (1) of Theorem 17.4. The goal of this problem is to compute the polynomials $u, v \in \mathbb{C}[x]$ for which condition (2) of Theorem 17.4 holds, and the endomorphism ϕ for which condition (3) of Theorem 17.4 holds, and to explicitly confirm that the diagram

$$\begin{array}{ccc} \mathbb{C}/L & \xrightarrow{\quad} & E(\mathbb{C}) \\ | & & | \\ \downarrow & & \downarrow \\ \mathbb{C}/L & \xrightarrow{\quad \phi \quad} & E(\mathbb{C}) \end{array}$$

commutes, where τ denotes the multiplication-by- τ map $z \mapsto \tau z$.

Recall that the Weierstrass \wp -function satisfying the differential equation

$$\wp'(z)^2 = 4\wp(z)^3 - g_2\wp(z) - g_3 \tag{1}$$

has a Laurent series expansion about 0 of the form $\wp(z) = z^{-2} + \sum_{n=1}^{\infty} a_{2n}z^{2n}$.

- (a) Use g_2 and g_3 to determine a_2 and a_4 , and then determine a_6 by comparing coefficients in the Laurent expansions of both sides of (1).

We now wish to compute the polynomials $u, v \in \mathbb{C}[x]$ for which

$$\wp(\tau z) = \frac{u(\wp(z))}{v(\wp(z))},$$

as in condition (2) of Theorem 17.4. Following Corollary 17.5, we have $N(\tau) = \tau\bar{\tau} = 2$, so $\deg u = 2$ and $\deg v = 1$. We can make $u = x^2 + ax + b$ monic, and with $v = cx + d$ we must have

$$(c\wp(z) + d)\wp(\tau z) = \wp(z)^2 + a\wp(z) + b \quad (2)$$

- (b) Use (2) to determine the coefficients a, b, c, d , expressing your answers in terms of τ . It will be convenient to work in the subfield $K = \mathbb{Q}(\tau)$, rather than \mathbb{C} . To define the field K and the polynomial ring $K[x]$ in Sage, use

```
RQ.<w>=PolynomialRing(QQ)
K.<tau>=NumberField(w^2-w+2)
RK.<x>=PolynomialRing(K)
```

Once you have determined $a, b, c, d \in K$, you can verify $u, v \in K[x]$ via¹

```
RL.<z>=LaurentSeriesRing(K, 100)
wp=EllipticCurve([-35, -98]).weierstrass_p(100).change_ring(K)
assert wp(tau*z) == u(wp(z))/v(wp(z))
```

- (c) Following the proof of Theorem 17.4, construct polynomials $s, t \in K[x]$ that satisfy

$$\wp'(\tau z) = \frac{s(\wp(z))}{t(\wp(z))} \wp'(z).$$

You can verify your results in Sage via

```
wpp = wp.derivative()
assert wpp(tau*z) == s(wp(z))/t(wp(z))*wpp(z)
```

- (d) Now let $\phi = \left(\frac{u(x)}{v(x)}, \frac{s(x)}{t(x)}y\right)$. Use Sage to verify that ϕ is an endomorphism by checking that its coordinate functions satisfy the curve equation $y^2 = 4x^3 - g_2x - g_3$.

The symbolic verifications in parts (b) and (d) confirm that $\Phi(\tau z) = \phi(\Phi(z))$, showing that the diagram commutes (at least for the first 100 terms in the Laurent expansion of $\wp(z)$). But we would like to explicitly check this for some specific values of $z \in \mathbb{C}$. In order to do this in Sage, we need to redefine τ and the polynomials u, v, s, t over \mathbb{C} , rather than K . You can use the following Sage script to do this:

```
R.<X>=PolynomialRing(CC)
pi = K.embeddings(CC)[0]
tauC = pi(tau)
def coerce(f, pi, X):
    c = f.coefficients(sparse=False)
    return sum([pi(c[i])*X^i for i in range(len(c))])
uC = coerce(u, pi, X)
vC = coerce(v, pi, X)
sC = coerce(s, pi, X)
tC = coerce(t, pi, X)
```

¹Sage effectively computes $\wp(z)$ using $y^2 = 4x^3 - g_2x - g_3$ when we define $E: y^2 = x^3 + Ax + B$ with $g_2 = -4A$ and $g_3 = -4B$.

- (e) Pick three “random” nonzero complex numbers z_1, z_2, z_3 of norm less than 0.1 (they need to be close to 0 in order for the Laurent series of $\wp(x)$ to converge quickly). You can approximate the point $P_1 = \Phi(z_1) = (\wp(z_1), \wp'(z_1))$ on the elliptic curve $y^2 = 4x^3 - g_2x - g_3$ in Sage using²

```
wp = EllipticCurve([CC(-35), CC(-98)]).weierstrass_p(100)
wpp = wp.derivative()
P1=(wp.laurent_polynomial()(z1), wpp.laurent_polynomial()(z1))
```

For $i = 1, 2, 3$, compute the points $P_i = \Phi(z_i)$ and $Q_i = \Phi(\tau z_i)$ (remember to use the embedding of τ in \mathbb{C}). Check that the points all approximately satisfy the curve equation $y^2 = 4x^3 - g_2x - g_3$ (if not, use z_i with smaller norms). Then verify that Q_i and $\phi(P_i)$ are approximately equal in each case. Report the values of z_i, P_i, Q_i and $\phi(P_i)$.

Problem 2. Binary quadratic forms (49 points)

A *binary quadratic form* is a homogeneous polynomial of degree 2 in two variables:

$$f(x, y) = ax^2 + bxy + cy^2,$$

which we identify by the coefficient vector (a, b, c) . We are interested in a particular set of binary quadratic forms, those that are *integral* ($a, b, c \in \mathbb{Z}$), *primitive* ($\gcd(a, b, c) = 1$), and *positive definite* ($b^2 - 4ac < 0$ and $a > 0$). Henceforth we shall use the word *form* to refer to an integral, primitive, positive definite, binary quadratic form. The *discriminant* of a form is the negative integer $D = b^2 - 4ac$, which is evidently a square modulo 4. We call such integers (imaginary quadratic) discriminants, and let $F(D)$ denote the set of forms with discriminant D .

- (a) For each $\gamma = \begin{pmatrix} s & t \\ u & v \end{pmatrix} \in \mathrm{SL}_2(\mathbb{Z})$ and $f(x, y) \in F(D)$ define

$$f^\gamma(x, y) := f(sx + ty, ux + vy).$$

Show that $f^\gamma \in F(D)$, and that this defines a right group action of $\mathrm{SL}_2(\mathbb{Z})$ on the set $F(D)$ (this means $f^I = f$ and $f^{(\gamma_1\gamma_2)} = (f^{\gamma_1})^{\gamma_2}$).

Forms f and g are (properly) *equivalent* if $g = f^\gamma$ for some $\gamma \in \mathrm{SL}_2(\mathbb{Z})$. In this problem and the next, you will prove that the set $\mathrm{cl}(D)$ of $\mathrm{SL}_2(\mathbb{Z})$ -equivalence classes of $F(D)$ forms a finite abelian group, and develop algorithms to compute in this group.

The group $\mathrm{cl}(D)$ is called the *class group*, and it plays a key role in the theory of complex multiplication. Our first objective is to prove that $\mathrm{cl}(D)$ is finite, and to develop an algorithm to enumerate unique representatives of its elements (which also allows us to determine its cardinality). We define the (principal) *root* τ of a form $f = (a, b, c)$ to be the unique root of $f(x, 1)$ in the upper half plane:

$$\tau = \frac{-b + \sqrt{D}}{2a}.$$

Recall that $\mathrm{SL}_2(\mathbb{Z})$ acts on the upper half plane \mathbb{H} via linear fractional transformations

$$\begin{pmatrix} s & t \\ u & v \end{pmatrix} \tau = \frac{s\tau + t}{u\tau + v},$$

²You need to use the `laurent_polynomial` method in order to evaluate `wp` at a complex number.

and that the set

$$\mathcal{F} = \{\tau \in \mathbb{H} : \operatorname{re}(\tau) \in [-1/2, 0] \text{ and } |\tau| \geq 1\} \cup \{\tau \in \mathbb{H} : \operatorname{re}(\tau) \in (0, 1/2) \text{ and } |\tau| > 1\}$$

is a fundamental region for \mathbb{H} modulo the $\mathrm{SL}_2(\mathbb{Z})$ -action.

- (b) Prove that $\gamma \in \mathrm{SL}_2(\mathbb{Z})$ acts compatibly on forms and their roots by showing that if τ is the root of f , then $\gamma^{-1}\tau$ is the root of f^γ . Conclude that two forms are equivalent if and only if their roots are equivalent.

A form $f = (a, b, c)$ is said to be *reduced* if

$$-a < b \leq a < c \quad \text{or} \quad 0 \leq b \leq a = c.$$

- (c) Prove that a form is reduced if and only if its root lies in the fundamental region \mathcal{F} . Conclude that each equivalence class in $F(D)$ contains exactly one reduced form.
- (d) Prove that if f is reduced then $a \leq \sqrt{|D|/3}$. Conclude that the set $\mathrm{cl}(D)$ is finite, and show that in fact its cardinality $h(D)$ satisfies $h(D) \leq |D|/3$. Prove that $F(D)$ contains a unique reduced form (a, b, c) with $a = 1$, and conclude that $h(-3) = h(-4) = 1$.

The positive integer $h(D)$ is called the *class number* of the discriminant D . The bound $h(D) \leq |D|/3$ is a substantial overestimate. In fact, $h(D) = O(|D|^{1/2} \log |D|)$, but proving this requires some analytic number theory that is beyond the scope of this course. Under the generalized Riemann hypothesis one can show $h(D) = O(|D|^{1/2} \log \log |D|)$.

- (e) Give an algorithm to enumerate the reduced forms in $F(D)$. Using the upper bound $h(D) = O(|D|^{1/2} \log |D|)$, prove that your algorithm runs in $O(|D|M(\log |D|))$ time.
- (f) Implement your algorithm and use it to enumerate the five reduced forms in $F(-103)$ and the six reduced forms in $F(-396)$. Then use it to compute $h(D)$ for the first three discriminants $D < -N$, where N is the integer formed by the first four digits of your student ID.

Problem 3. The class group (98 points)

In Problem 2 we proved that $\mathrm{cl}(D)$ is a finite set. In this problem you will prove that it is an abelian group, and develop an algorithm for computing the group operation.

To each form $f(x, y) = ax^2 + bxy + cy^2$ in $F(D)$ with root $\tau = (-b + \sqrt{D})/(2a)$, we associate the lattice $L(f) = L(a, b, c) = a[1, \tau]$.

- (a) Show that two forms $f, g \in F(D)$ are equivalent if and only if the lattices $L(f)$ and $L(g)$ are homothetic (use may use part (b) of problem 2 if you wish).

For any lattice L , the *order* of L is the set

$$\mathcal{O}(L) = \{\alpha \in \mathbb{C} : \alpha L \subseteq L\}.$$

- (b) Prove that either $\mathcal{O}(L) = \mathbb{Z}$ or $\mathcal{O}(L)$ is an order in an imaginary quadratic field, and that homothetic lattices have the same order. Prove that if L is the lattice of a form in $F(D)$, then $\mathcal{O}(L)$ is the order of discriminant D in the field $K = \mathbb{Q}(\sqrt{D})$.

For the rest of this problem let \mathcal{O} denote the (not necessarily maximal) imaginary quadratic order of discriminant D , which may be represented as a lattice $[1, \omega]$, where ω is an algebraic integer whose minimal polynomial $x^2 + bx + c$ has discriminant $b^2 - 4c = D$.

Recall that an (integral) \mathcal{O} -ideal \mathfrak{a} is an additive subgroup of \mathcal{O} that is closed under multiplication by \mathcal{O} . Every \mathcal{O} -ideal \mathfrak{a} is necessarily a sublattice of \mathcal{O} , and its *norm* $N(\mathfrak{a})$ is the index $[\mathcal{O} : \mathfrak{a}] = |\mathcal{O}/\mathfrak{a}|$. An \mathcal{O} -ideal \mathfrak{a} is said to be *proper* if $\mathcal{O}(\mathfrak{a}) = \mathcal{O}$. In Lecture 18 we showed that \mathfrak{a} is proper if and only if it is invertible as a fractional ideal, which explains our interest in this property. Note that we always have $\mathcal{O} \subseteq \mathcal{O}(\mathfrak{a})$, so when \mathcal{O} is maximal every nonzero \mathcal{O} -ideal is proper.

- (c) Prove that if $L(a, b, c) = a[1, \tau]$ is the lattice of a form in $F(D)$, then L is a proper \mathcal{O} -ideal of norm a , where $\mathcal{O} = \mathcal{O}(L) = [1, a\tau]$.
- (d) Conversely prove that every proper \mathcal{O} -ideal \mathfrak{a} is homothetic to the lattice of a form in $F(D)$. Show that the assumption that \mathfrak{a} is proper is necessary by giving an explicit example of an \mathcal{O} -ideal \mathfrak{a} that is not proper (so by (c) it cannot be homothetic to the lattice of a form in $F(D)$).
- (e) Prove that if the norm of \mathfrak{a} is relatively prime to the conductor $u = [\mathcal{O}_K : \mathcal{O}]$ of \mathcal{O} then \mathfrak{a} is proper. Give an explicit example showing that the converse is not true.

The product of two lattices $[\omega_1, \omega_2]$ and $[\omega_3, \omega_4]$ in \mathbb{C} is the additive group generated by $\{\omega_1\omega_3, \omega_1\omega_4, \omega_2\omega_3, \omega_2\omega_4\}$.

- (f) Show that, in general, the product of two lattices need not be a lattice, but the product of two lattices that are \mathcal{O} -ideals is a lattice.
- (g) Let $\text{cl}(\mathcal{O})$ denote the set of equivalence classes (under homothety) of lattices that are proper \mathcal{O} -ideals. Prove that the lattice product makes $\text{cl}(\mathcal{O})$ into an abelian group. Conclude that the corresponding operation on the equivalence classes of $F(D)$ makes $\text{cl}(D)$ into an abelian group that is isomorphic to $\text{cl}(\mathcal{O})$.

To do explicit computations in $\text{cl}(D)$ we need to translate the product operation on lattices $L(f_1)$ and $L(f_2)$ into a corresponding product operation on forms $f_1, f_2 \in F(D)$. This is known as *composition* of forms, and is performed as follows. If $f_1 = (a_1, b_1, c_1)$ and $f_2 = (a_2, b_2, c_2)$ are forms in $F(D)$, then let $s = (b_1 + b_2)/2$ (this is an integer because b_1, b_2 and D all have the same parity). Use the extended Euclidean algorithm (twice) to compute integers u, v, w , and d such that $ua_1 + va_2 + ws = d = \gcd(a_1, a_2, s)$. The composition of f_1 and f_2 is then given by

$$f_1 * f_2 = (a_3, b_3, c_3) = \left(\frac{a_1 a_2}{d^2}, b_2 + \frac{2a_2}{d}(v(s - b_2) - wc_2), \frac{b_3^2 - D}{4a_3} \right).$$

It is a straight-forward but tedious task to verify that this composition formula satisfies $L(f_1 * f_2) = L(f_1) * L(f_2)$; you are not required to do this.

- (h) Verify that the inverse of (a, b, c) is $(a, -b, c)$ and that the unique reduced form with $a = 1$ acts as the identity (see Problem 2 for the definition of a reduced form).

Unfortunately, even if f_1 and f_2 are reduced forms, the composition of f_1 and f_2 need not be reduced. In order to compute in $\text{cl}(D)$ effectively, we need a reduction algorithm. Recall the matrices $S = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$ and $T = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ that generate $\text{SL}_2(\mathbb{Z})$.

- (i) Let f be the form (a, b, c) . Compute the forms f^S , f^{T^m} , and $f^{T^{-m}}$, for a positive integer m .

A form (a, b, c) with $-a < b \leq a$ is said to be *normalized*.

- (j) Show that for any form f there is an integer m such that f^{T^m} is normalized, and give an explicit formula for m . Let us call f^{T^m} the *normalization* of f . Now let $f = (a, b, c)$ be a normalized form and prove the following:

- (a) If $a < \sqrt{|D|}/2$ then f is reduced.
 (b) If $a < \sqrt{|D|}$ and f is not reduced, then the normalization of Sf is reduced.
 (c) If $a \geq \sqrt{|D|}$ then the normalization (a', b', c') of Sf has $a' \leq a/2$.

- (k) Give an algorithm to compute the reduction of a form f in $F(D)$, and bound its complexity as a function of $n = \log |D|$, assuming that its coefficients are $O(n)$ bits in size. Then bound the complexity of computing the reduction of the product of two reduced forms (this corresponds to performing a group operation in $\text{cl}(D)$).³

- (l) Implement your algorithm and then use it to compute the reduction of a form $(a, b, c) \in F(D)$, with a equal to the least prime greater than $|D|^2$ for which $(\frac{D}{a}) = 1$. Do this for the discriminants $D = -103$ and $D = -396$, and for the first three discriminants $D < -N$, where N is the first four digits of your student ID. For the largest $|D|$, list the sequence of normalized forms computed during the reduction.

Problem 4. Subgroups of $\text{GL}_2(\mathbb{F}_\ell)$ (49 points)

Let E be an elliptic curve defined over \mathbb{Q} . Recall that for each integer $n > 1$, the n -torsion subgroup of $E(\overline{\mathbb{Q}})$ is a rank 2 $(\mathbb{Z}/n\mathbb{Z})$ -module we denote $E[n]$. As explained in Problem Sets 3 and 6, the action of the absolute Galois group $G_{\mathbb{Q}} := \text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q})$ on the coordinates of points gives rise to an action on the set $E(\overline{\mathbb{Q}})$ that commutes with the group law. Hence the $G_{\mathbb{Q}}$ -action preserves $E[n]$ and gives rise to a linear representation of the absolute Galois group

$$\rho_{E,n}: G_{\mathbb{Q}} \rightarrow \text{Aut}(E[n]) \simeq \text{GL}_2(\mathbb{Z}/n\mathbb{Z}),$$

which we call the *mod- n Galois representation* attached to E . In this problem we restrict our attention to prime $n = \ell$, in which case we have following theorem of Serre.

Theorem (Serre, 1972). *Let E be an elliptic curve over \mathbb{Q} for which $\text{End}(E_{\overline{\mathbb{Q}}}) = \mathbb{Z}$. For all but finitely many primes ℓ , the image of the mod- ℓ Galois representation is surjective:*

$$\rho_{E,\ell}(G_{\mathbb{Q}}) = \text{GL}_2(\mathbb{F}_\ell).$$

Remark. For an elliptic curve over \mathbb{Q} (or any number field) we know that $\text{End}(E_{\overline{\mathbb{Q}}})$ is either \mathbb{Z} or an order in an imaginary quadratic field. The latter case is quite special: it applies to only 13 \mathbb{Q} -isomorphism classes of elliptic curves over \mathbb{Q} , corresponding to the 13 imaginary quadratic orders of class number one.⁴

³A quasi-linear bound is known [1], but your bound does not need to be this tight. However it should be polynomial in n .

⁴Elliptic curves E/\mathbb{Q} with $\text{End}(E_{\overline{\mathbb{Q}}}) \neq \mathbb{Z}$ are often said to have complex multiplication and called CM curves, even though this is not strictly true: the extra endomorphisms are only defined over a quadratic extension (it would be more correct to say these curves have “potential complex multiplication”).

Remark. It is conjectured that Serre's theorem actually applies to all primes $\ell > 37$ (independent of E). There is ample evidence and some recent progress toward a proof of this conjecture, but it remains a major open question.

A key component of the proof of Serre's theorem is understanding the maximal subgroups of $\mathrm{GL}_2(\mathbb{F}_\ell)$. In order to discuss subgroups of $\mathrm{GL}_2(\mathbb{F}_\ell)$ in a basis-free manner, it is often convenient to write $\mathrm{GL}(V)$ where V is a 2-dimensional vector space over \mathbb{F}_ℓ and $\mathrm{GL}(V)$ denotes its group of automorphisms. In this problem you will give a complete classification of the maximal subgroups of $\mathrm{GL}_2(V)$.

Let L_1 and L_2 be distinct 1-dimensional subspaces of V , which we can think of as lines through the origin in V , and let C_s be the subgroup of $\mathrm{GL}(V)$ that preserves both L_1 and L_2 (individually, no swapping allowed).

- (a) Show that for $\ell \neq 2$, the subgroup C_s uniquely determines the lines $L_1, L_2 \subset V$ (and hence is equivalent to specifying two such lines).

We call such a C a *split Cartan subgroup* of $\mathrm{GL}(V)$. If we choose a basis for V compatible with the decomposition $V = L_1 \oplus L_2$, we then have

$$C_s = \begin{pmatrix} * & 0 \\ 0 & * \end{pmatrix},$$

where $*$ indicates any element of \mathbb{F}_ℓ^\times . From this we see that $C \simeq (\mathbb{F}_\ell^\times)^2$ is an abelian group of order $(\ell - 1)^2$.

As an \mathbb{F}_ℓ -vector space, $\mathbb{F}_{\ell^2} \simeq \mathbb{F}_\ell^2$; but \mathbb{F}_{ℓ^2} also has a multiplicative structure, and so the action of the multiplicative group $\mathbb{F}_{\ell^2}^\times$ on $\mathbb{F}_{\ell^2} \simeq V$ gives a cyclic subgroup C_{ns} of $\mathrm{GL}(V)$ isomorphic to $\mathbb{F}_{\ell^2}^\times$. Such a subgroup C_{ns} is called a *non-split Cartan subgroup*. We collectively refer to split and non-split Cartan subgroups as Cartan subgroups.

- (b) Show that for $\ell \neq 2$, if we fix a quadratic non-residue $\epsilon \in \mathbb{F}_\ell^\times$, then in an appropriate basis we have

$$C_{ns} = \left\{ \begin{pmatrix} x & \epsilon y \\ y & x \end{pmatrix} : x, y \in \mathbb{F}_\ell, (x, y) \neq (0, 0) \right\}.$$

- (c) Show that the intersection of any two distinct Cartan subgroups (either split or non-split) is the group of scalar matrices $Z = \begin{pmatrix} z & 0 \\ 0 & z \end{pmatrix}$ with $z \in \mathbb{F}_\ell^\times$.
- (d) Show that any element $s \in \mathrm{GL}(V)$ with $\Delta(s) = \mathrm{tr}(s)^2 - 4 \cdot \det(s) \neq 0$ is contained in a unique Cartan subgroup, and determine a condition involving $\Delta(s)$ that specifies the type of Cartan. Deduce that the union of all Cartan subgroups of $\mathrm{GL}(V)$ is the set of elements of order prime to ℓ . (If you are stuck, look at part (h) below.)
- (e) Let N denote the normalizer of a Cartan subgroup C in $\mathrm{GL}(V)$, that is all elements $s \in \mathrm{GL}(V)$ such that $sCs^{-1} = C$. Show that $(N : C) = 2$ and give an explicit description of this group in the split and non-split cases separately.

It is easy to show that the group Z of scalar matrices forms the center of $\mathrm{GL}(V)$. We define $\mathrm{PGL}(V)$ to be the quotient of $\mathrm{GL}(V)$ by its center, so $\mathrm{PGL}(V) := \mathrm{GL}(V)/Z$. Let $\varphi: \mathrm{GL}(V) \rightarrow \mathrm{PGL}(V)$ denote the quotient map.

- (f) Show that if C is a split (resp. non-split) Cartan subgroup, then $\varphi(C) \subset \mathrm{PGL}(V)$ is cyclic of order $\ell - 1$ (resp. $\ell + 1$). Show that the image in $\mathrm{PGL}(V)$ of a normalizer of a Cartan subgroup is a dihedral group.⁵

By part (d) above, it remains to understand the elements of $\mathrm{GL}(V)$ of order divisible by ℓ . A Borel subgroup B of $\mathrm{GL}(V)$ is the group of automorphisms of V fixing a specified line (through the origin). A Borel subgroup of $\mathrm{GL}(V)$ has order $\ell(\ell - 1)^2$. After choosing an appropriate basis, this has the form

$$B = \begin{pmatrix} * & * \\ 0 & * \end{pmatrix}.$$

- (g) Show that any element $s \in \mathrm{GL}(V)$ of order ℓ is conjugate to the matrix $\begin{pmatrix} 1 & \\ 0 & 1 \end{pmatrix}$.
- (h) Using the fact that $\mathrm{SL}(V)$ is generated $\begin{pmatrix} 1 & \\ 0 & 1 \end{pmatrix}$ and $\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$, deduce that any subgroup of $\mathrm{GL}(V)$ of order divisible by ℓ either lies in a Borel subgroup, or contains $\mathrm{SL}(V)$.

Let k be any field. If H is a finite subgroup of $\mathrm{PGL}_2(k)$ of order prime to the characteristic of k that is not cyclic or dihedral, then H is isomorphic to either A_4, S_4 , or A_5 . (In the case $k = \mathbb{C}$, this result is well known; these subgroups correspond to the symmetry groups of the regular polyhedra: tetrahedron, cube/octahedron, and icosahedron/dodecahedron, respectively.)

- (i) Use parts (a) to (h) to prove the following classification theorem.

Theorem (Maximal subgroups of $\mathrm{GL}_2(\mathbb{F}_\ell)$). *Let G be a subgroup of $\mathrm{GL}_2(\mathbb{F}_\ell)$; let H denote the image of G in $\mathrm{PGL}_2(\mathbb{F}_\ell)$. Then one of the following holds:*

1. G has order prime to ℓ and either:
 - (i) H is cyclic and G is contained in a Cartan subgroup of $\mathrm{GL}_2(\mathbb{F}_\ell)$;
 - (ii) H is dihedral and G is contained in the normalizer of a Cartan subgroup C of $\mathrm{GL}_2(\mathbb{F}_\ell)$ but not in C ;
 - (iii) H is isomorphic to A_4, S_4 or A_5 and we call G exceptional;
2. G has order divisible by ℓ and either:
 - (iv) G is contained in a Borel subgroup;
 - (v) G contains $\mathrm{SL}_2(\mathbb{F}_\ell)$.

Serre's theorem states that except for elliptic curves E/\mathbb{Q} with (potential) complex multiplication, for all but finitely many primes ℓ we are in case (v) of the classification above. On later problem sets we will see that for $\ell \neq 2$ this never happens if E has complex multiplication, so the hypothesis $\mathrm{End}(E_{\overline{\mathbb{Q}}}) = \mathbb{Z}$ in Serre's theorem is necessary.

⁵For this problem, the product of two cyclic groups of order 2 (the Klein group) is a dihedral group.

Problem 5. Survey (2 points)

Complete the following survey by rating each of the problems you attempted on a scale of 1 to 10 according to how interesting you found the problem (1 = “mind-numbing,” 10 = “mind-blowing”), and how difficult you found it (1 = “trivial,” 10 = “brutal”). Also estimate the amount of time you spent on each problem to the nearest half hour.

| | Interest | Difficulty | Time Spent |
|-----------|----------|------------|------------|
| Problem 1 | | | |
| Problem 2 | | | |
| Problem 3 | | | |
| Problem 4 | | | |

Also, please rate each of the following lectures that you attended, according to the quality of the material (1=“useless”, 10=“fascinating”), the quality of the presentation (1=“epic fail”, 10=“perfection”), the pace (1=“way too slow”, 10=“way too fast”, 5=“just right”) and the novelty of the material (1=“old hat”, 10=“all new”).

| Date | Lecture Topic | Material | Presentation | Pace | Novelty |
|------|------------------------|----------|--------------|------|---------|
| 4/10 | Complex multiplication | | | | |
| 4/17 | The CM torsor | | | | |

Please feel free to record any additional comments you have on the problem sets or lectures, in particular, ways in which they might be improved.

References

- [1] A. Schönhage, *Fast reduction and composition of binary quadratic forms*, in International Symposium on Symbolic and Algebraic Computation–ISSAC’91, ACM, 1991, 128–133.

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