

MIT 6.035
Semantic Analysis

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Error Issue

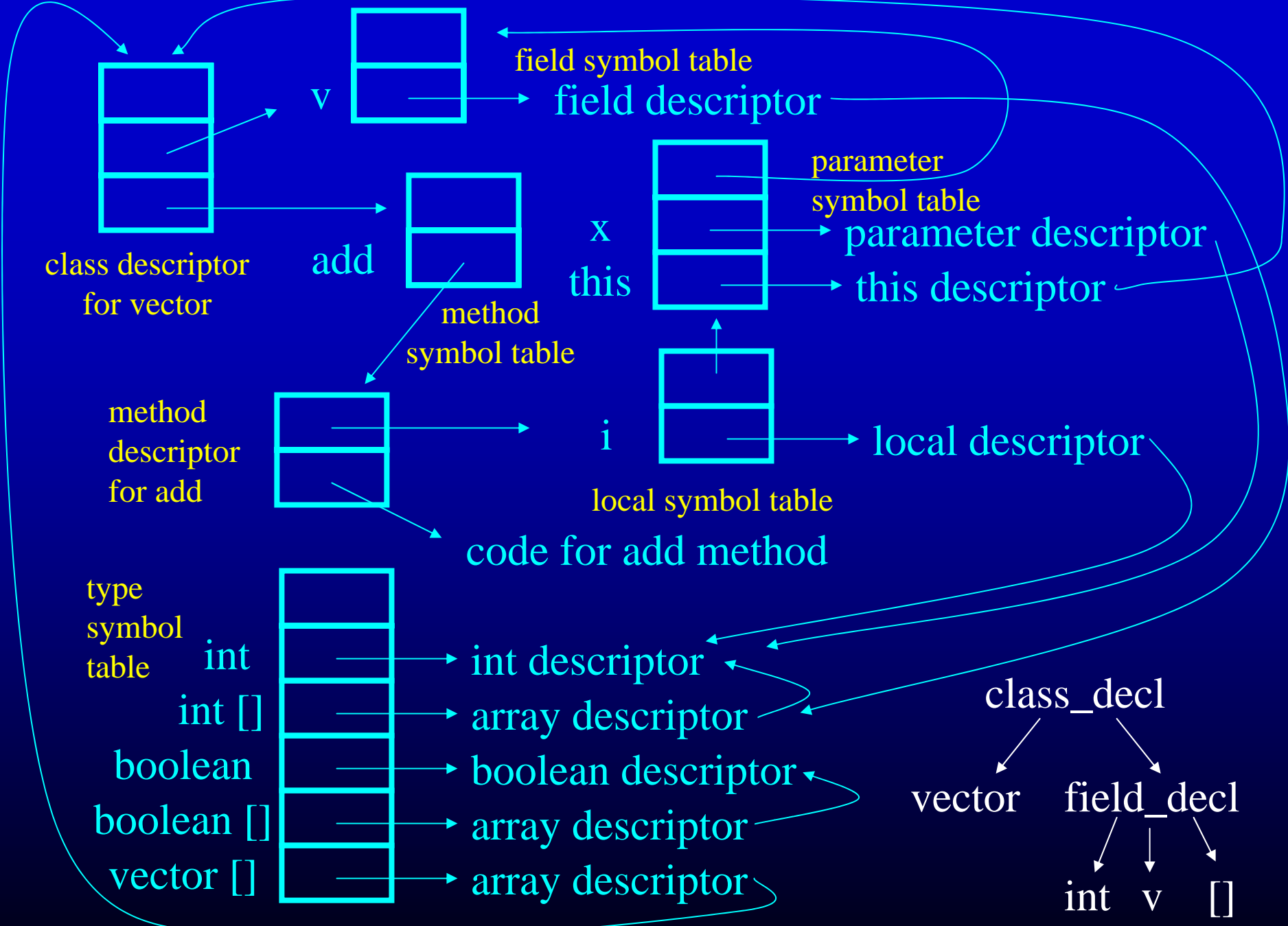
- Have assumed no problems in building IR
- But are many static checks that need to be done as part of translation
- Called Semantic Analysis

Goal of Semantic Analysis

- Ensure that program obeys certain kinds of sanity checks
 - all used variables are defined
 - types are used correctly
 - method calls have correct number and types of parameters and return value
- Checked when build IR
- Driven by symbol tables

Symbol Table Summary

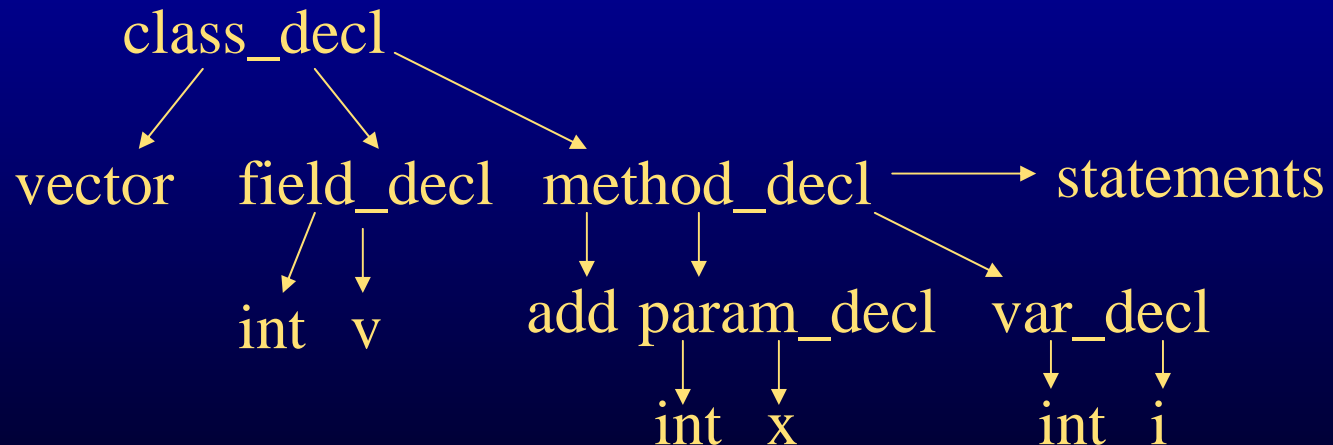
- Program Symbol Table (Class Descriptors)
- Class Descriptors
 - Field Symbol Table (Field Descriptors)
 - Field Symbol Table for SuperClass
 - Method Symbol Table (Method Descriptors)
 - Method Symbol Table for Superclass
- Method Descriptors
 - Local Variable Symbol Table (Local Variable Descriptors)
 - Parameter Symbol Table (Parameter Descriptors)
 - Field Symbol Table of Receiver Class
- Local, Parameter and Field Descriptors
 - Type Descriptors in Type Symbol Table or Class Descriptors

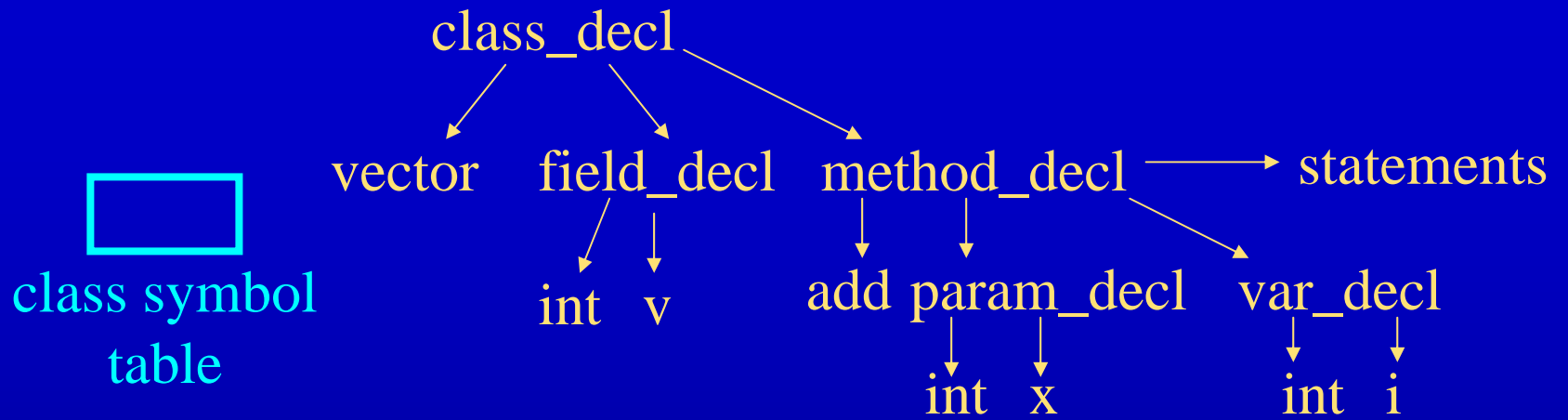


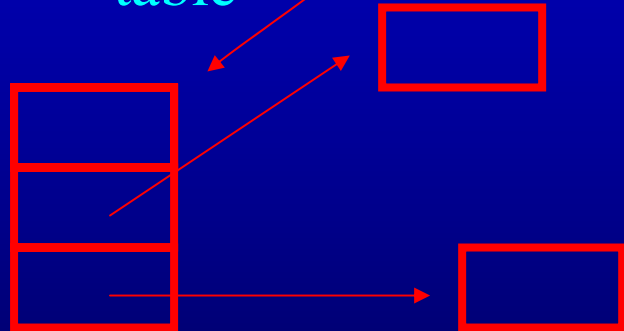
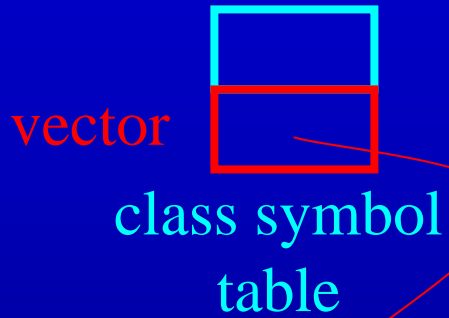
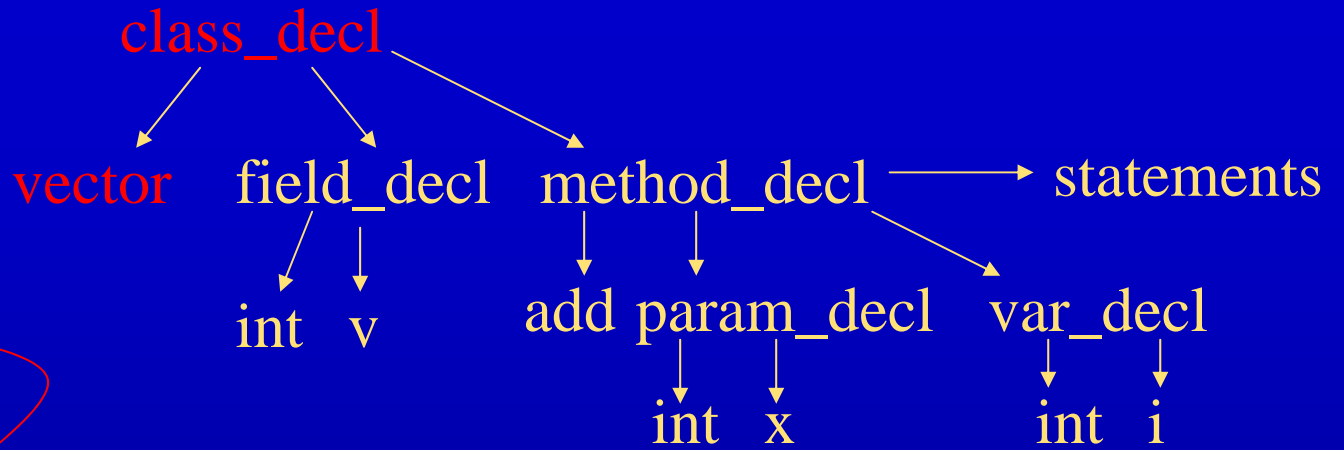
Translating from Abstract Syntax Trees to Symbol Tables

Intermediate Representation for Classes

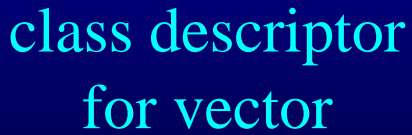
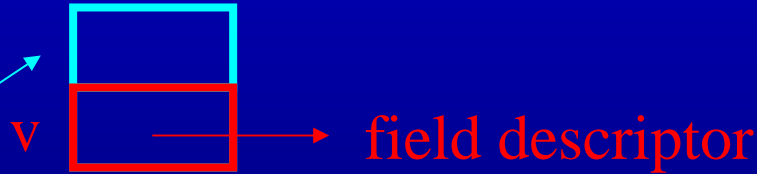
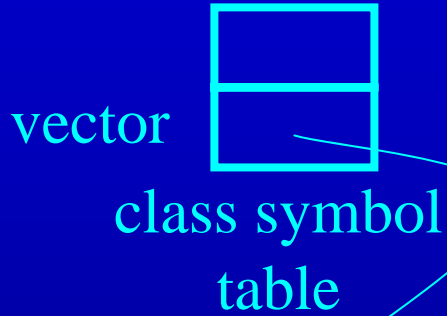
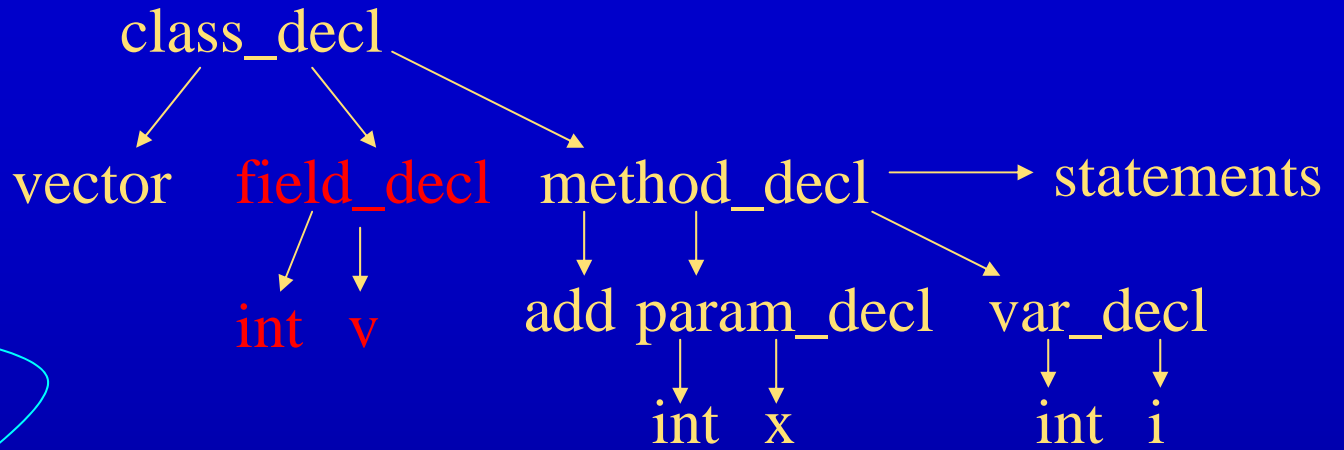
```
class vector {  
    int v[];  
    void add(int x) {  
        int i; i = 0;  
        while (i < v.length) { v[i] = v[i]+x; i = i+1; }  
    }  
}
```

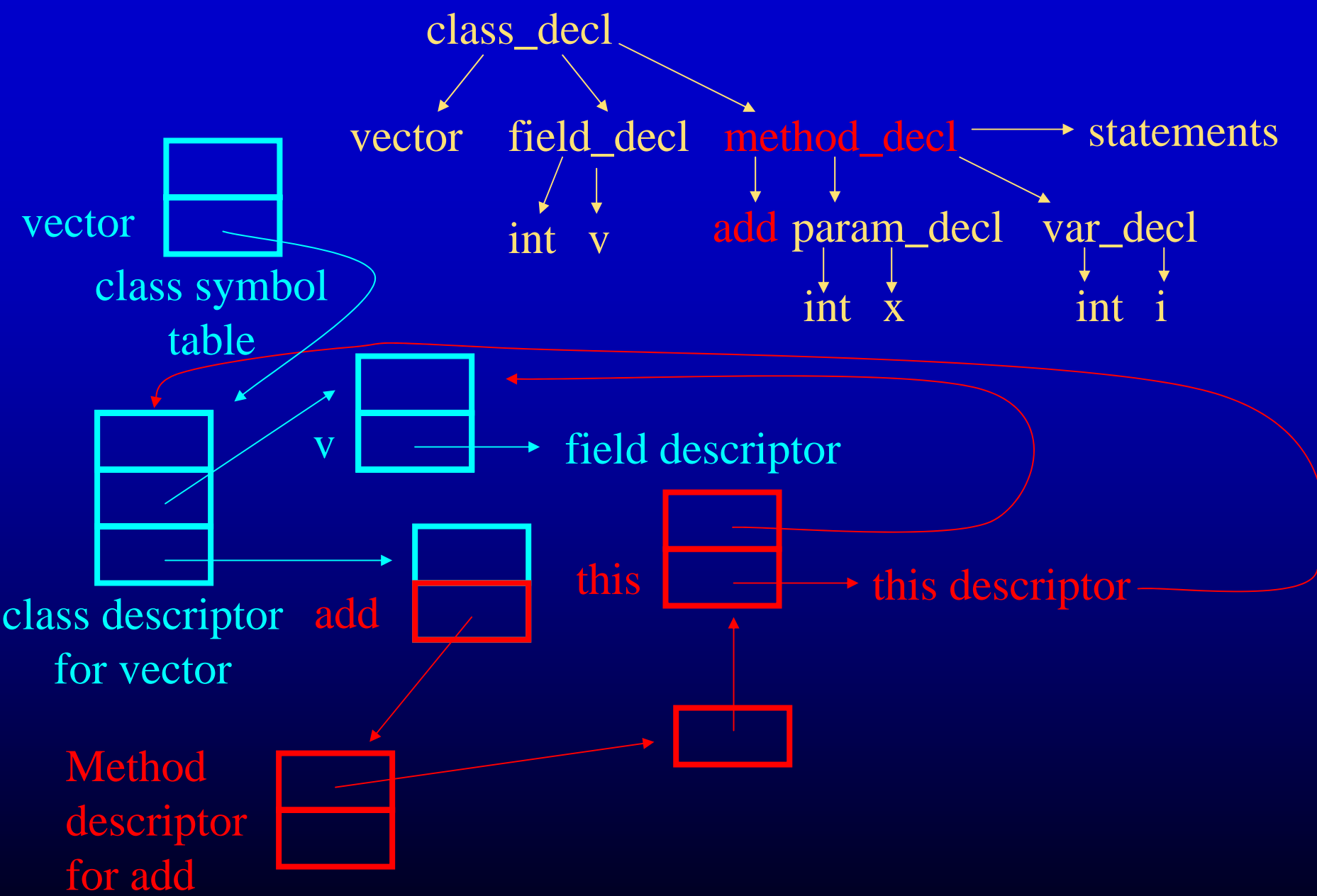


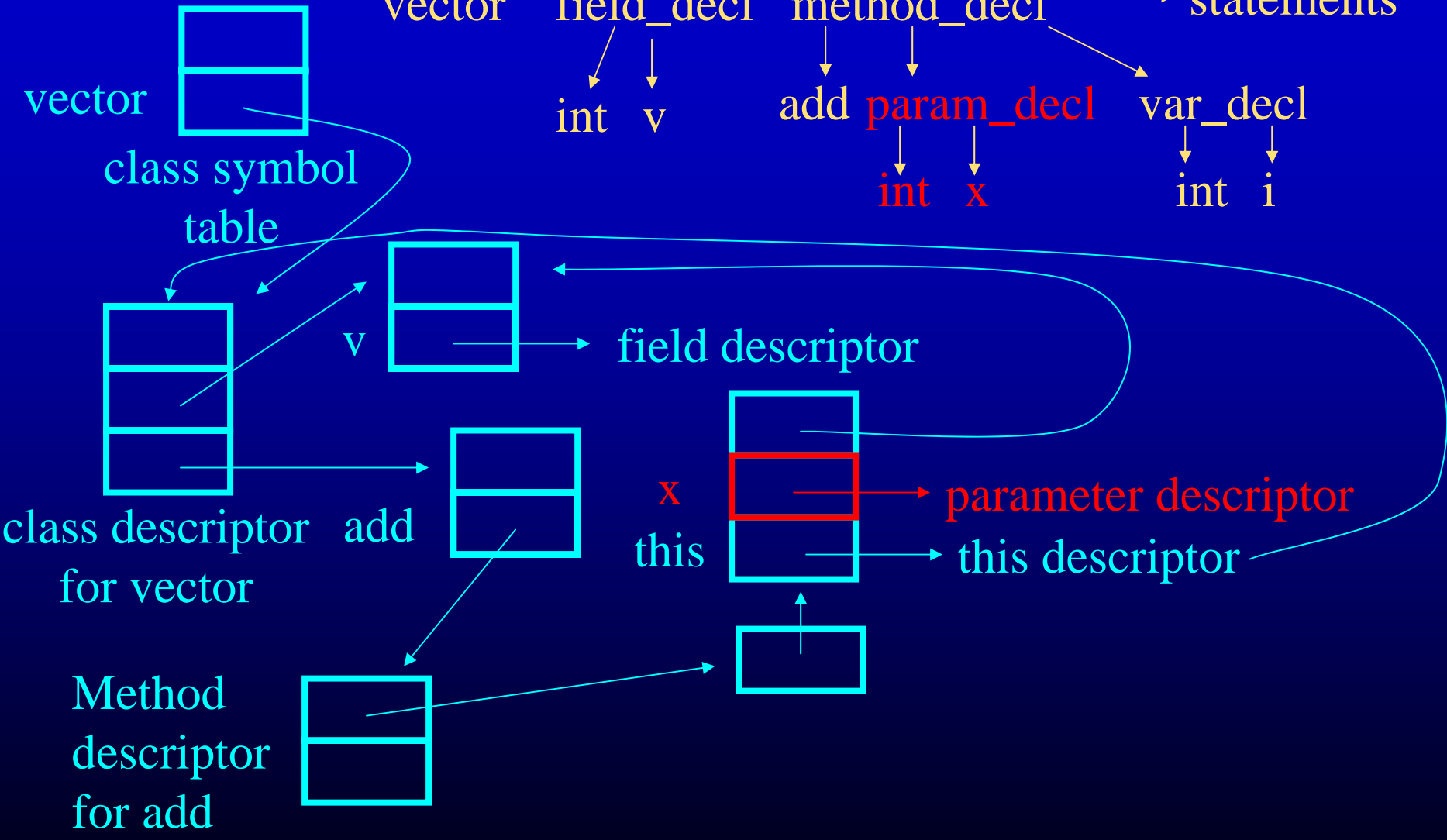
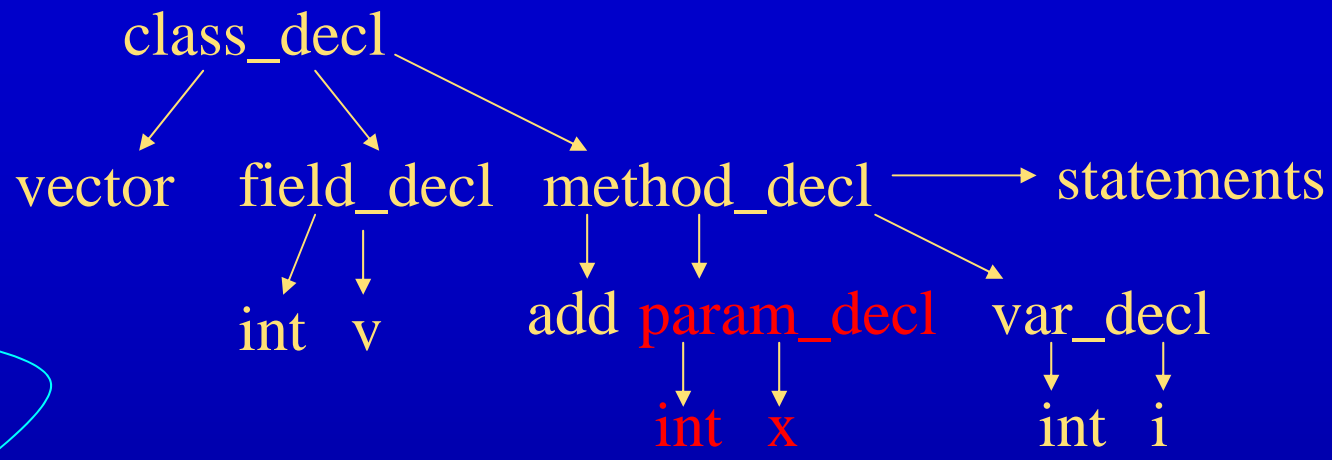


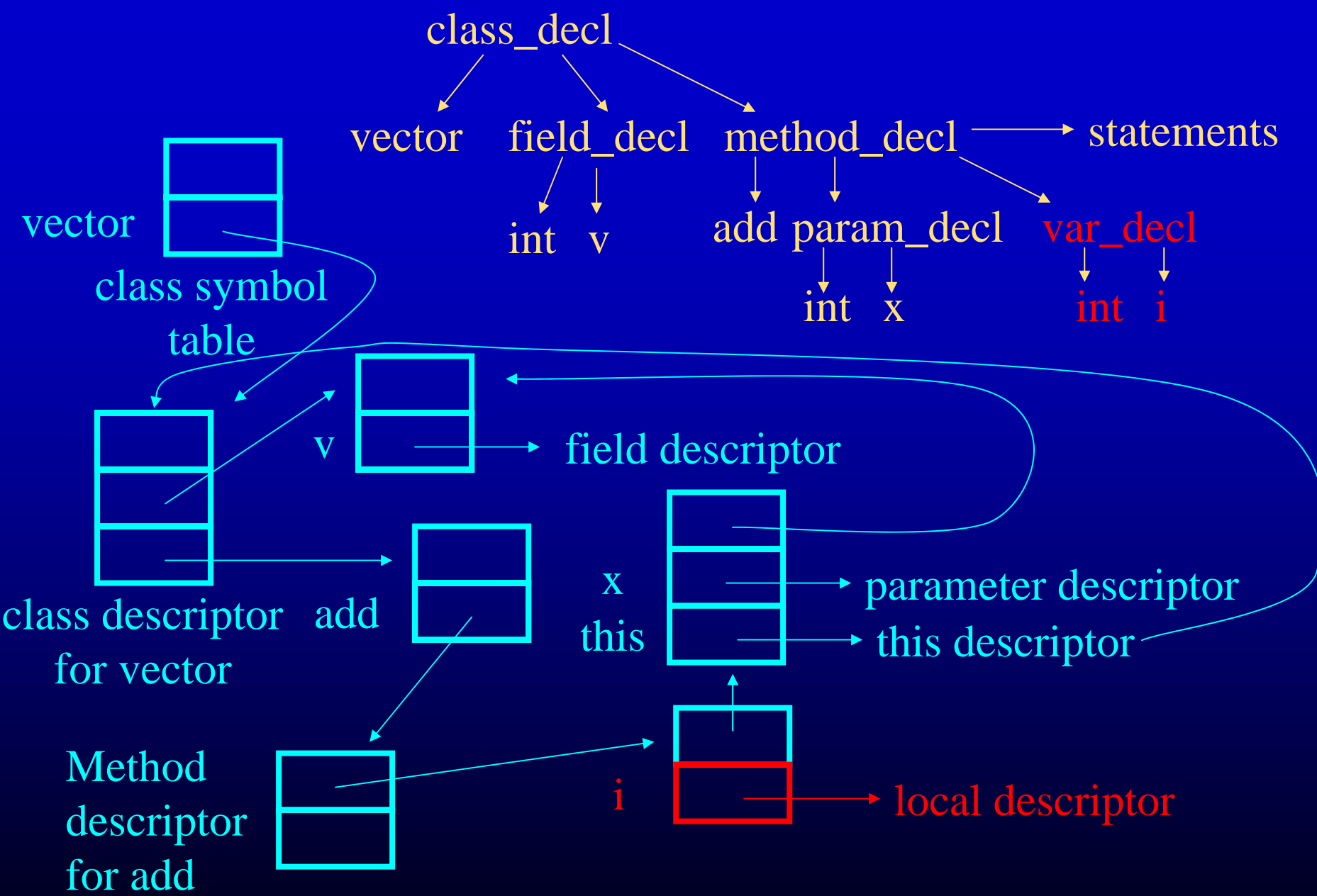


class descriptor for vector





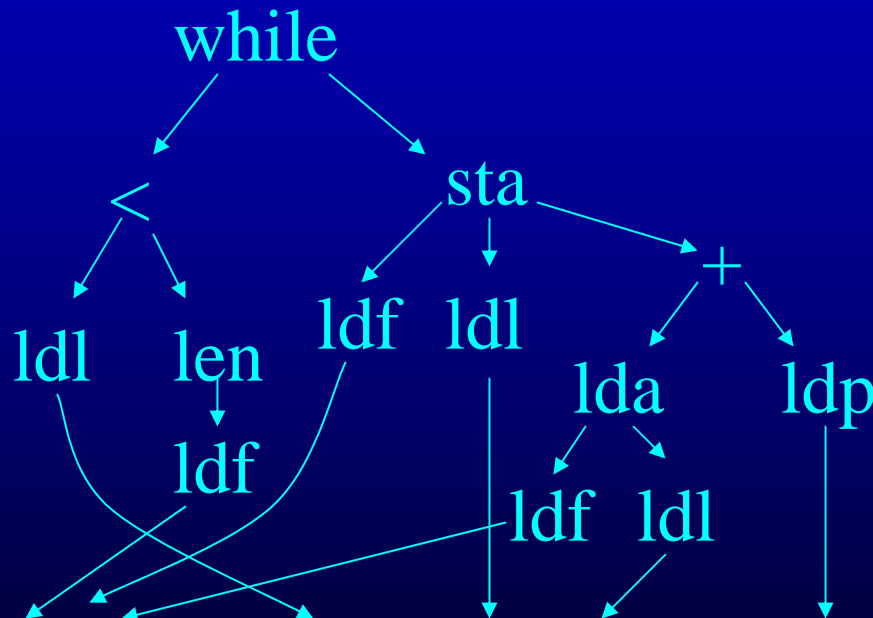




Intermediate Representation for Code

while (i < v.length)

v[i] = v[i]+x;



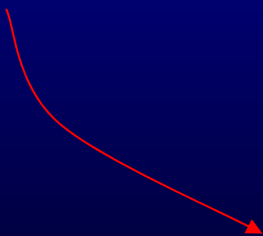
field descriptor for v local descriptor for i parameter descriptor for x

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    v[i] = v[i]+x;
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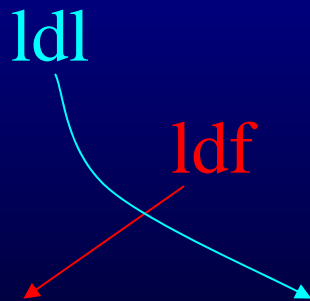
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```

ldl



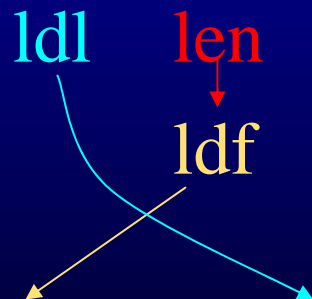
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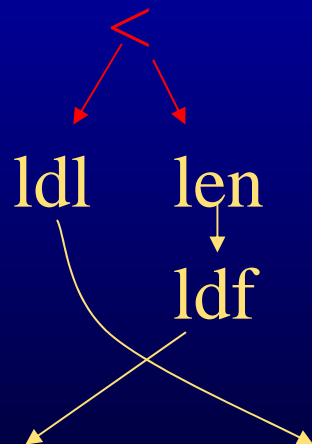
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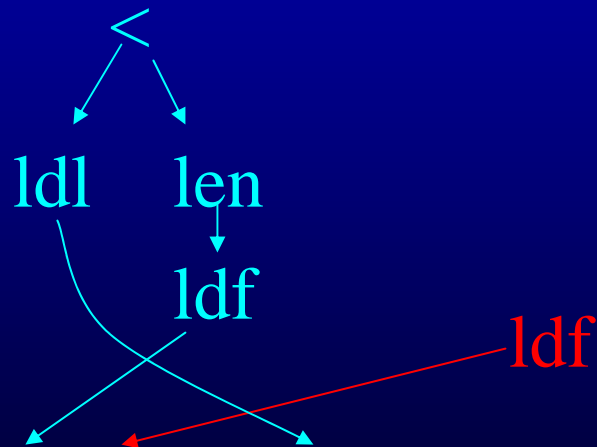
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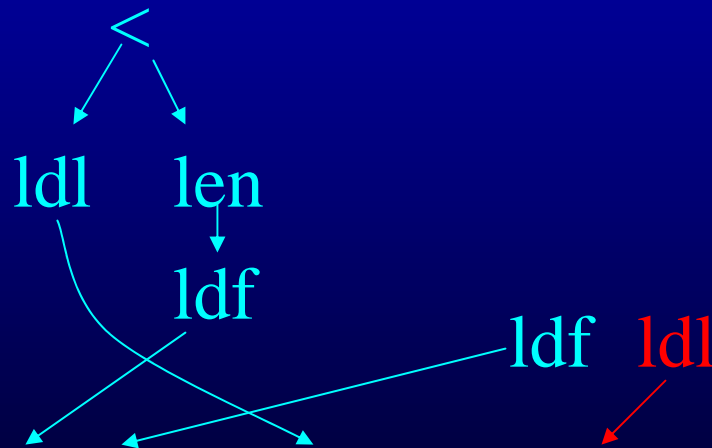
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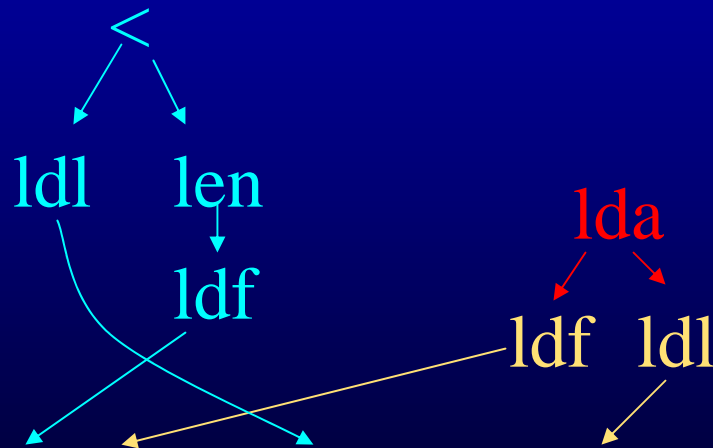
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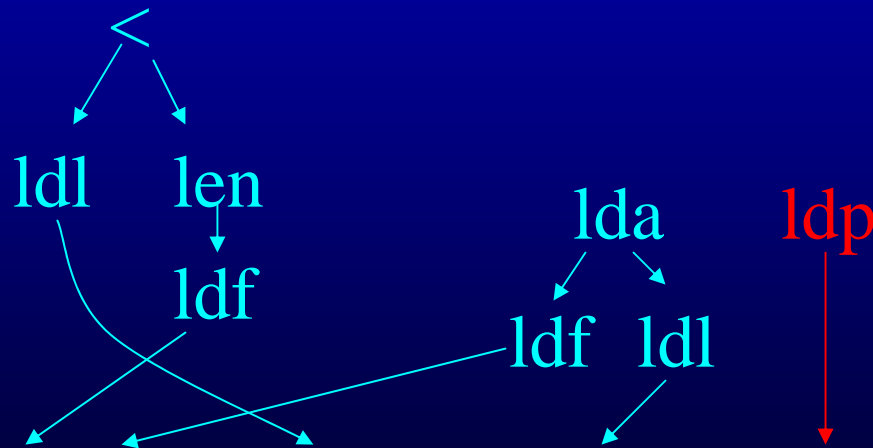
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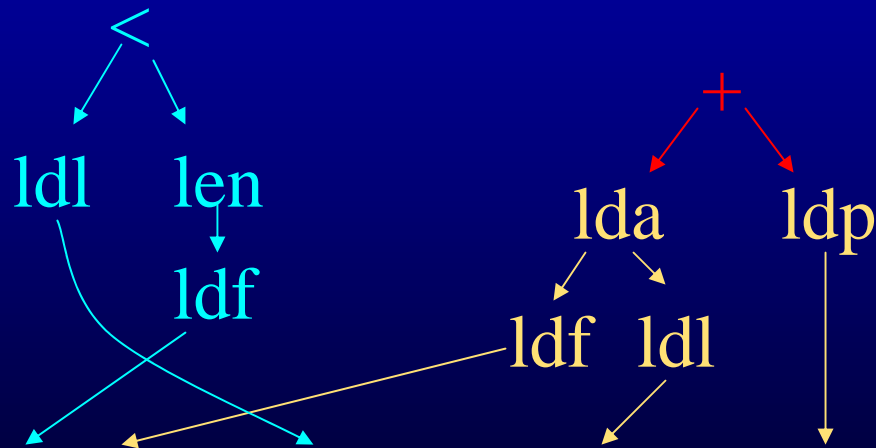
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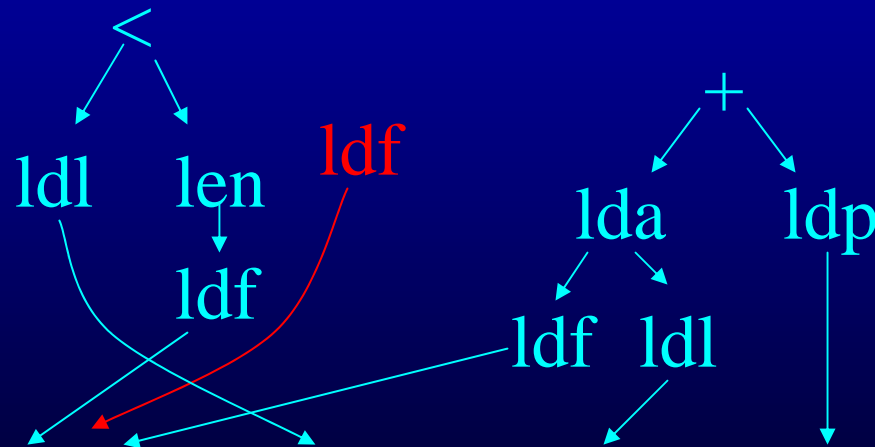
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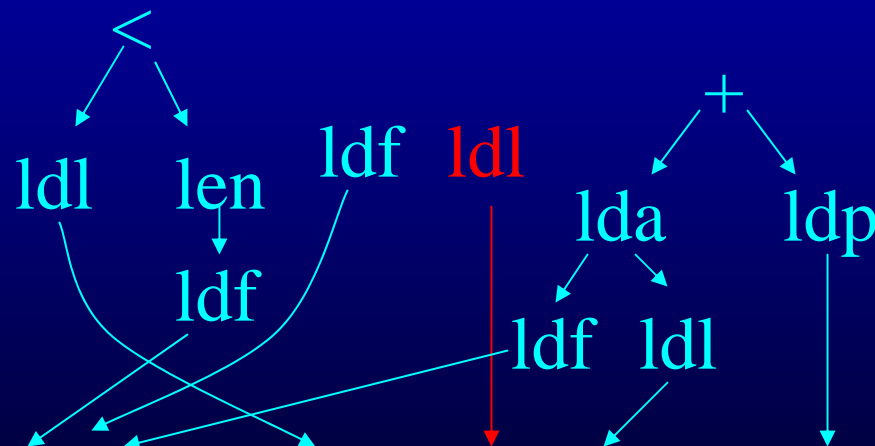
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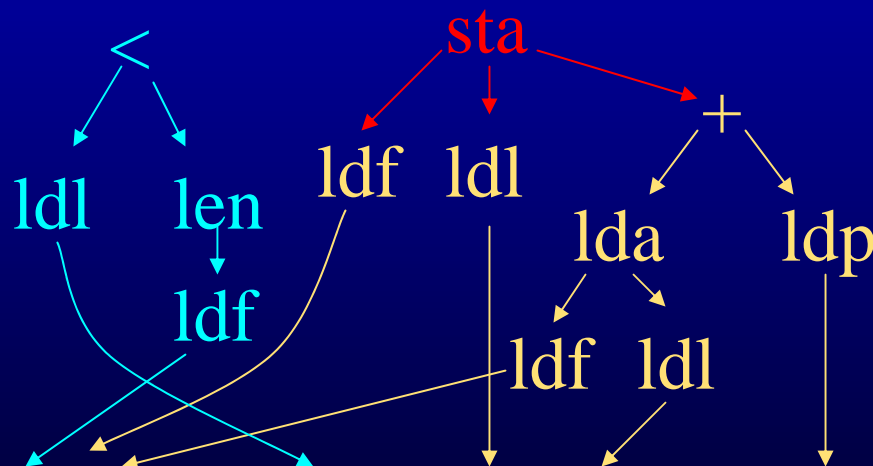
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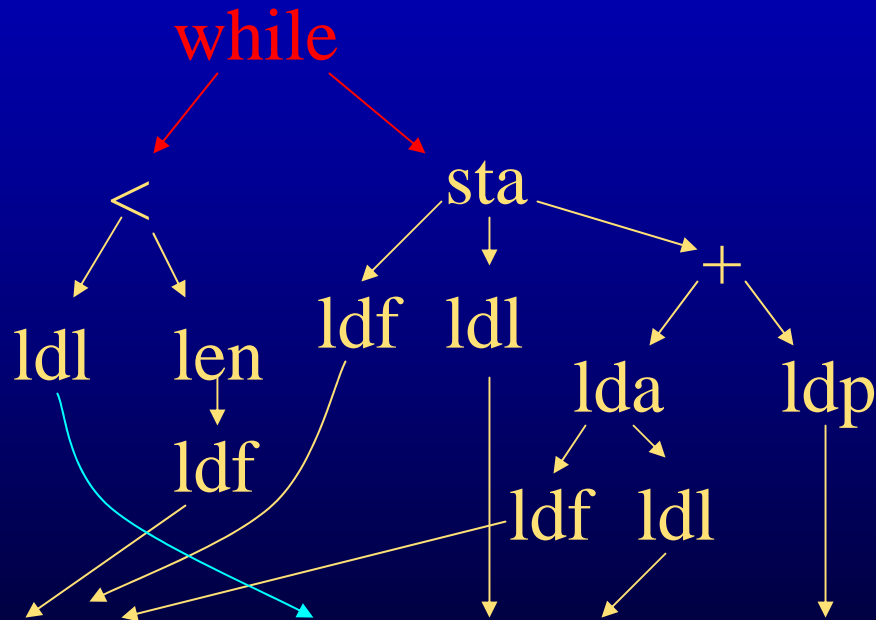
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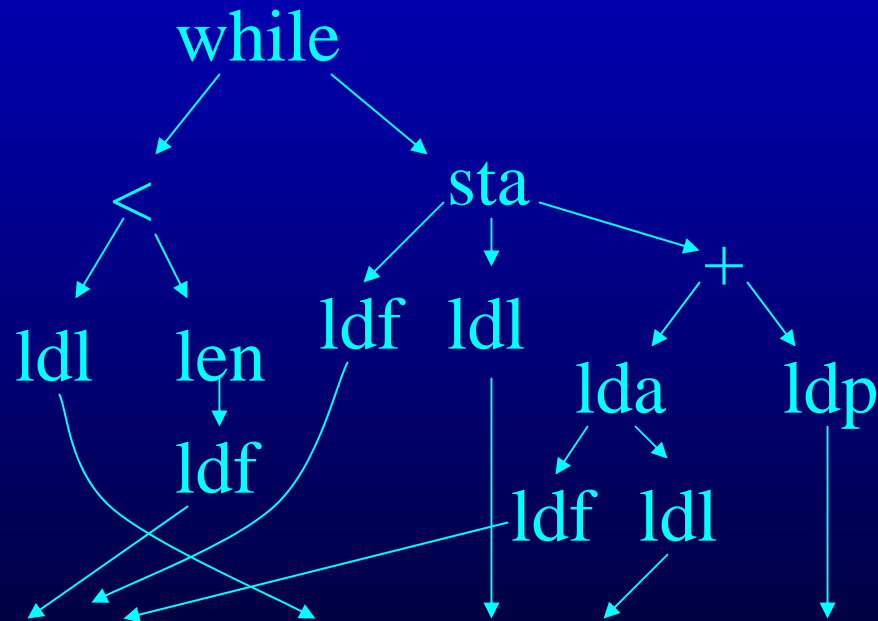
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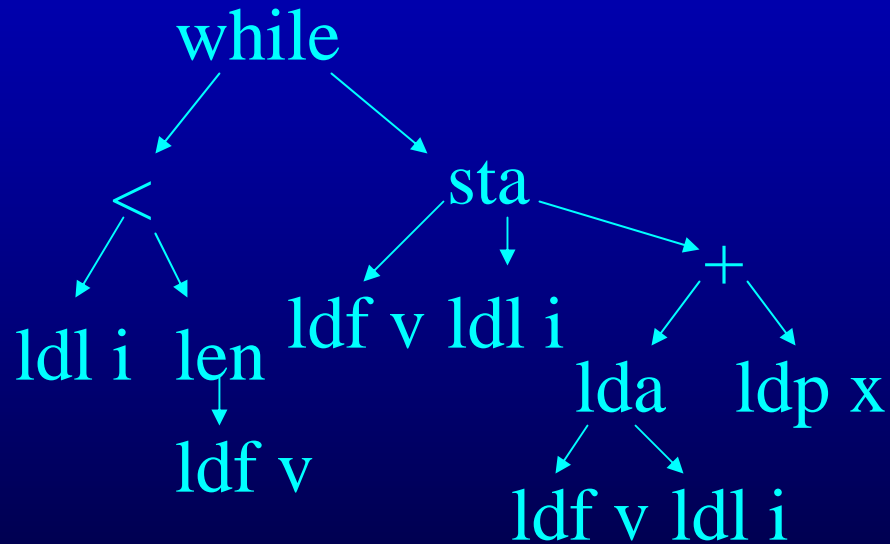


field descriptor for v local descriptor for i parameter descriptor for x

Abbreviated Notation

while (i < v.length)

v[i] = v[i]+x;



Parameter Descriptors

- When build parameter descriptor, have
 - name of type
 - name of parameter
- What is the check? Must make sure name of type identifies a valid type
 - look up name in type symbol table
 - if not there, look up name in program symbol table (might be a class type)
 - if not there, fails semantic check

Local Descriptors

- When build local descriptor, have
 - name of type
 - name of local
- What is the check? Must make sure name of type identifies a valid type
 - look up name in type symbol table
 - if not there, look up name in program symbol table (might be a class type)
 - if not there, fails semantic check

Local Symbol Table

- When build local symbol table, have a list of local descriptors
- What to check for?
 - duplicate variable names
 - shadowed variable names
- When to check?
 - when insert descriptor into local symbol table
- Parameter and field symbol tables similar

Class Descriptor

- When build class descriptor, have
 - class name and name of superclass
 - field symbol table
 - method symbol table
- What to check?
 - Superclass name corresponds to actual class
 - No name clashes between field names of subclass and superclasses
 - Overridden methods match parameters and return type declarations of superclass

Load Instruction

- What does compiler have? Variable name.
- What does it do? Look up variable name.
 - If in local symbol table, reference local descriptor
 - If in parameter symbol table, reference parameter descriptor
 - If in field symbol table, reference field descriptor
 - If not found, semantic error

Load Array Instruction

- What does compiler have?
 - Variable name
 - Array index expression
- What does compiler do?
 - Look up variable name (if not there, semantic error)
 - Check type of expression (if not integer, semantic error)

Add Operations

- What does compiler have?
 - two expressions
- What can go wrong?
 - expressions have wrong type
 - must both be integers (for example)
- So compiler checks type of expressions
 - load instructions record type of accessed variable
 - operations record type of produced expression
 - so just check types, if wrong, semantic error

Type Inference for Add Operations

- Most languages let you add floats, ints, doubles
- What are issues?
 - Types of result of add operation
 - Coercions on operands of add operation
- Standard rules usually apply
 - If add an int and a float, coerce the int to a float, do the add with the floats, and the result is a float.
 - If add a float and a double, coerce the float to a double, do the add with the doubles, result is double

Add Rules

- Basic Principle: Hierarchy of number types (int, then float, then double)
- All coercions go up hierarchy
 - int to float; int, float to double
- Result is type of operand highest up in hierarchy
 - int + float is float, int + double is double, float + double is double
- Interesting oddity: C converts float procedure arguments to doubles. Why?

Type Inference

- Infer types without explicit type declarations
- Add is very restricted case of type inference
- Big topic in recent programming language research
 - How many type declarations can you omit?
 - Tied to polymorphism

Equality Expressions

- If build expression $A = B$, must check compatibility
- A compatible with B or B compatible with A
- Int compatible with Int
- Class C compatible with Class D if
C inherits from D (but not vice-versa)

Store Instruction

- What does compiler have?
 - Variable name
 - Expression
- What does it do?
 - Look up variable name.
 - If in local symbol table, reference local descriptor
 - If in parameter symbol table, error
 - If in field symbol table, reference field descriptor
 - If not found, semantic error
 - Check type of variable name against type of expression
 - If variable type not compatible with expression type, error

Store Array Instruction

- What does compiler have?
 - Variable name, array index expression
 - Expression
- What does it do?
 - Look up variable name.
 - If in local symbol table, reference local descriptor
 - If in parameter symbol table, error
 - If in field symbol table, reference field descriptor
 - If not found, semantic error
 - Check that type of array index expression is integer
 - Check type of variable name against type of expression
 - If variable element type not compatible with expression type, error

Method Invocations

- What does compiler have?
 - method name, receiver expression, actual parameters
- Checks:
 - receiver expression is class type
 - method name is defined in receiver's class type
 - types of actual parameters match types of formal parameters
 - What does match mean?
 - same type?
 - compatible type?

Semantic Check Summary

- Do semantic checks when build IR
- Many correspond to making sure entities are there to build correct IR
- Others correspond to simple sanity checks
- Each language has a list that must be checked
- Can flag many potential errors at compile time