

Lecture 5

Rationalizability

14.12 Game Theory

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Recall: A Game

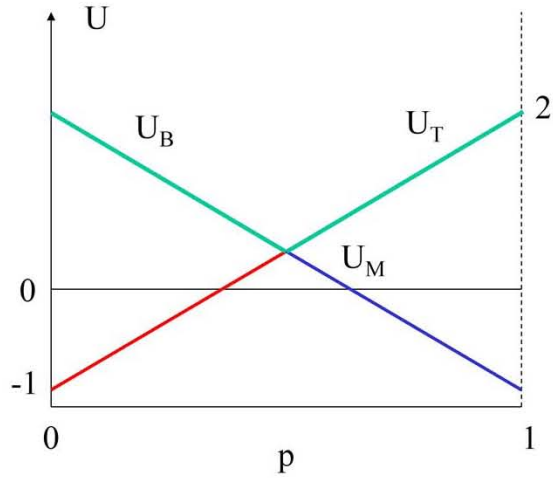
$$U_M = 0$$

$$U_T = 2p - (1-p) = 3p - 1$$

$$U_B = -p + 2(1-p) = 2 - 3p$$

	L	R
T	(2,0)	(-1,1)
M	(0,10)	(0,0)
B	(-1,-6)	(2,0)

p 1-p



Recap: Rationality & Dominance

- **Belief**: A probability distribution p_{-i} on **others'** strategies;
- **Mixed Strategy**: A probability distribution σ_i on **own** strategies;
- Playing s_i^* is **rational** $\Leftrightarrow s_i^*$ is a **best response** to a belief p_{-i} : $\forall s_{-i}$
$$\sum_{s_{-i}} u_i(s_i^*, s_{-i}) p_{-i}(s_{-i}) \geq \sum_{s_{-i}} u_i(s_i, s_{-i}) p_{-i}(s_{-i})$$
- σ_i **dominates** s_i^{**} $\Leftrightarrow \forall s_{-i}$
$$\sum_{s_{-i}} u_i(s_i, s_{-i}) \sigma_i(s_i) > u_i(s_i^{**}, s_{-i})$$
- Theorem: Playing s_i^* is rational $\Leftrightarrow s_i^*$ is not dominated.

Assume

Player 1 is rational

Player 2 is rational

Player 2 is rational and

Knows that Player 1 is rational

Player 1 is rational,

knows that 2 is rational

knows that 2 knows that

1 is rational

	L	R
T	(2,0)	(-1,1)
M	(0,10)	(0,0)
B	(-1,-6)	(2,0)

Assume

P1 is rational

P2 is rational and
knows that P1 is rational

	2	L	m	R
1				
T		(3,0)	(1,1)	(0,3)
M		(1,0)	(0,10)	(1,0)
B		(0,3)	(1,1)	(3,0)

P1 is rational and
knows all these

Rationalizability



The play is rationalizable, provided that ...

Important

- Eliminate only the **strictly** dominated strategies
 - Ignore weak dominance
- Make sure to eliminate the strategies dominated by mixed strategies as well as pure

Beauty Contest

- There are n students.
- Simultaneously, each student submits a number x_i between 0 and 100.
- The payoff of student i is $100 - (x_i - 2\bar{x}/3)^2$ where

$$\bar{x} = \frac{x_1 + x_2 + \dots + x_n}{n}.$$

Rationalizability in Beauty Contest

If X_{-i} = Expected value of sum of x_j with $j \neq i$, best strategy is
 $(2/3)X_{-i}/(n-2/3)$

After Round 1:

$$\left[0, \frac{2}{3} \frac{n-1}{n-2/3} 100 \right]$$

After Round 2:

$$\left[0, \left(\frac{2}{3} \frac{n-1}{n-2/3} \right)^2 100 \right]$$

After Round k :

$$\left[0, \left(\frac{2}{3} \frac{n-1}{n-2/3} \right)^k 100 \right]$$

Rationalizability = $\{0\}$.

with m mischievous students

Payoff for mischievous: $(x_i - 2x/3)^2$

Round 1: only 0 and 100 survive for mischievous; same as before for normal

Rounds 2 to $k(m,n)-1$: no elimination for mischievous; same as before for normal

Round $k(m,n)$: eliminate 0 for mischievous; same as before for normal

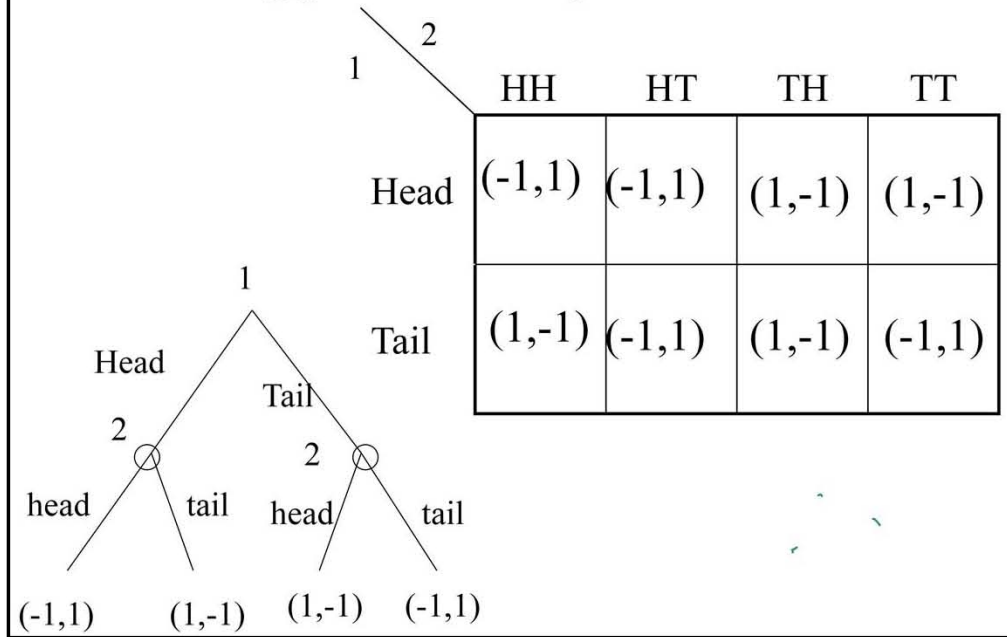
Round $k > k(m,n)$:

– Strategies for normal after round $k = [L_k, H_k]$

$$L_k = \frac{2}{3} \frac{100m + (n-m-1)L_{k-1}}{n-2/3} \quad H_k = \frac{2}{3} \frac{100m + (n-m-1)H_{k-1}}{n-2/3}$$

Rationalizability = mischievous 100, normal $200m/(n+2m)$

Matching pennies with perfect information



A summary

- If players are rational and cautious, they play the dominant-strategy equilibrium whenever it exists
 - But, typically, it does not exist
- If rationality is common knowledge, a rationalizable strategy is played
 - Typically, there are too many rationalizable strategies
- **Nash Equilibrium**: the players correctly guess the other players' strategies (or conjectures).

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14.12 Economic Applications of Game Theory
Fall 2012

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